
8 The Device Kit

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8 The Device Kit

The Device Kit contains software for controlling the various kinds of connectors and devices that can be attached to a BeBox and for writing your own drivers for input/output devices. You'll find two kinds of software documented in this chapter:

- Encapsulated interfaces to some of the ports found on the back of the BeBox. Currently, this part of the kit contains just one class—BSerialPort. A BSerialPort object can represent any of the four RS-232 serial ports that are visible on the back of the machine. Other classes for other connectors will be added in future releases.
- The programming interface and protocols for developing device drivers. All drivers are dynamically loaded, add-on modules that run as extensions either of the Application Server (drivers for graphics cards) or of the kernel (all other drivers).

The BSerialPort class is part of the Be operating system library. The programming interface for device drivers is not. A driver links only to its host module, not to a library.

If you're interested in the interface to a serial port, begin reading on page 5; you can skip everything after page 15. If you're interesting in writing a device driver, skip the first part of the chapter and begin on page 17.

BSerialPort

Derived from: public BObject
Declared in: <device/SerialPort.h>

Overview

A BSerialPort object represents an RS-232 serial port connection to the BeBox. There are four such ports on the back of the machine.

Through BSerialPort functions, you can read data received at a serial port and write data over the connection. You can also configure the connection—for example, set the number of data and stop bits, determine the rate at which data is sent and received, and select the type of flow control (hardware or software) that should be used.

To read and write data, a BSerialPort object must first open one of the serial ports by name. For example:

```
BSerialPort *connection = new BSerialPort;  
if ( connection->Open("serial2") > 0 ) {  
    . . .  
}
```

The BSerialPort object communicates with the driver for the port it has open. The driver maintains an input buffer of 1K bytes to collect incoming data and an output buffer half that size to hold outgoing data. When the object reads and writes data, it reads from and writes to these buffers.

Constructor and Destructor

BSerialPort()

BSerialPort(void)

Initializes the BSerialPort object to the following default values:

- No flow control (see `SetFlowControl()`)
- A data rate of 19,200 bits per second (see `SetDataRate()`)
- A serial unit with 8 bits of data, 1 stop bit, and no parity (see `SetDataBits()`)
- A timeout of one tenth of a second to read a target of one byte (see `Read()`)

The new object doesn't represent any particular serial port. After construction, it's necessary to open one of the ports by name.

See also: **Open()**

~BSerialPort()

virtual ~BSerialPort(void)

Makes sure the port is closed before the object is destroyed.

Member Functions

ClearInput(), ClearOutput()

void ClearInput(void)

void ClearOutput(void)

These functions empty the serial port driver's input and output buffers, so that their contents won't be read (by the **Read()** function) and won't be transmitted over the connection (after having been written by **Write()**).

The buffers are cleared automatically when a port is opened.

See also: **Read()**, **Write()**, **Open()**

Close() see **Open()**

DataBits() see **SetDataBits()**

DataRate() see **SetDataRate()**

FlowControl() see **SetFlowControl()**

IsCTS()

bool IsCTS(void)

Returns **TRUE** if the Clear to Send (CTS) pin is asserted, and **FALSE** if not.

IsDCD()

```
bool IsDCD(void)
```

Returns **TRUE** if the Data Carrier Detect (DCD) pin is asserted, and **FALSE** if not.

IsDSR()

```
bool IsDSR(void)
```

Returns **TRUE** if the Data Set Ready (DSR) pin is asserted, and **FALSE** if not.

IsRI()

```
bool IsRI(void)
```

Returns **TRUE** if the Ring Indicator (RI) pin is asserted, and **FALSE** if not.

Open(), Close()

```
long Open(const char *name)
```

```
void Close(void)
```

These functions open the *name* serial port and close it again. Ports are identified by names that correspond to their labels on the back panel of the BeBox:

```
“serial1”
```

```
“serial2”
```

```
“serial3”
```

```
“serial4”
```

To be able to read and write data, the BSerialPort object must have a port open. It can open first one port and then another, but it can have no more than one open at a time. If it already has a port open when **Open()** is called, that port is closed before an attempt is made to open the *name* port. (Thus, both **Open()** and **Close()** close the currently open port.)

Open() can't open the *name* port if some other entity already has it open. (If the BSerialPort itself has *name* open, **Open()** first closes it, then opens it again.)

If it's able to open the port, **Open()** returns a positive integer. If unable, it returns **B_ERROR**.

When a serial port is opened, its input and output buffers are emptied and the Data Terminal Ready (DTR) pin is asserted.

See also: **Read()**

ParityMode() see **SetDataBits()**

Read(), SetBlocking(), SetTimeout()

long **Read**(void **buffer*, long *maxBytes*)

void **SetBlocking**(bool *shouldBlock*)

void **SetTimeout**(double *timeout*)

Read() takes incoming data from the serial port driver and places it in the data *buffer* specified. In no case will it read more than *maxBytes*—a value that should reflect the capacity of the *buffer*; it returns the actual number of bytes read. **Read()** fails if the BSerialPort object doesn't have a port open.

The number of bytes that **Read()** will actually place in the buffer before returning depends on the *shouldBlock* flag and the *timeout* set by the other two functions.

SetBlocking() determines whether **Read()** should block and wait for *maxBytes* of data to arrive at the serial port if that number isn't already available to be read. If the *shouldBlock* flag is **TRUE**, **Read()** will block. However, if *shouldBlock* is **FALSE**, **Read()** will take however many bytes are waiting to be read, up to the maximum asked for, then return immediately. If no data is waiting at the serial port, it returns without reading anything.

The default *shouldBlock* setting is **TRUE**.

SetTimeout() sets a time limit on how long **Read()** will wait for a character to arrive at the input buffer. It's relevant only if the *shouldBlock* flag is **TRUE**. The *timeout* is expressed in microseconds and is limited to 255,000,000.0 (25.5 seconds); it's set to the maximum value if a greater amount of time is specified. Differences less than 100,000.0 microseconds (0.1 second) are not recognized; they're rounded to the nearest tenth of a second. There is no time limit if *timeout* is set to 0.0; **Read()** will block forever.

The default timeout is no timeout (0.0).

Like the standard **read()** system function, **Read()** returns the number of bytes it succeeded in placing in the *buffer*, which may be 0. It returns **SYS_ERROR** (−1) if there's an error on any kind—for example, if the BSerialPort object doesn't have a port open. It's not considered an error if a timeout expires.

See also: **Write()**, **Open()**

SetBlocking() see **Read()**

SetDataBits(), SetStopBits(), SetParityMode(), DataBits(), StopBits(), ParityMode()

```
void SetDataBits(data_bits count)
void SetStopBits(stop_bits count)
void SetParityMode(parity_mode mode)
data_bits DataBits(void)
stop_bits StopBits(void)
parity_mode ParityMode(void)
```

These functions set and return characteristics of the serial unit used to send and receive data. **SetDataBits()** sets the number of bits of data in each unit. The *count* can be:

```
B_DATA_BITS_7 or
B_DATA_BITS_8
```

The default is **B_DATA_BITS_8**.

SetStopBits() sets the number of stop bits in each unit. It can be:

```
B_STOP_BITS_1 or
B_STOP_BITS_2
```

The default is **B_STOP_BITS_1**.

SetParityMode() sets whether the serial unit contains a parity bit and, if so, the type of parity used. The mode can be:

```
B_EVEN_PARITY,
B_ODD_PARITY, or
B_NO_PARITY
```

The default is **B_NO_PARITY**.

SetDataRate(), DataRate()

```
void SetDataRate(data_rate bitsPerSecond)
data_rate DataRate(void)
```

These functions set and return the rate (in bits per second) at which data is both transmitted and received. Permitted values are:

B_0_BPS	B_200_BPS	B_4800_BPS
B_50_BPS	B_300_BPS	B_9600_BPS
B_75_BPS	B_600_BPS	B_19200_BPS
B_110_BPS	B_1200_BPS	B_38400_BPS
B_134_BPS	B_1800_BPS	B_57600_BPS
B_150_BPS	B_2400_BPS	B_115200_BPS

The default data rate is **B_19200_BPS**. If the rate is set to 0 (**B_0_BPS**), data will be sent and received at an indeterminate number of bits per second.

SetDTR()

long SetDTR(bool *pinAsserted*)

Asserts the Data Terminal Ready (DTR) pin if the *pinAsserted* flag is **TRUE**, and de-asserts it if the flag is **FALSE**. < This function isn't implemented for the current release. >

See also: **SetRTS()**

SetFlowControl(), FlowControl()

void SetFlowControl(ulong *mask*)

ulong FlowControl(void)

These functions set and return the type of flow control the driver should use. There are two possibilities:

B_SOFTWARE_CONTROL Control is maintained through XON and XOFF characters inserted into the data stream.

B_HARDWARE_CONTROL Control is maintained through the Clear to Send (CTS) and Request to Send (RTS) pins.

The *mask* passed to **SetFlowControl()** and returned by **FlowControl()** can be just one of these constants—or it can be a combination of the two, in which case the driver will use both types of flow control together. It can also be 0, in which case the driver won't use any flow control. **B_HARDWARE_CONTROL** is the default.

SetParityMode() see SetDataBits()

SetRTS()

long SetRTS(bool *pinAsserted*)

Asserts the Request to Send (RTS) pin if the *pinAsserted* flag is **TRUE**, and de-asserts it if the flag is **FALSE**. < This function isn't implemented for the current release. >

See also: **SetDTR()**

SetStopBits() see SetDataBits()

SetTimeout() see **Read()**

StopBits() see **SetDataBits()**

Write()

long **Write**(const void **data*, long *numBytes*)

Writes up to *numBytes* of *data* to the serial port's output buffer. This function will be successful in writing the data only if the BSerialPort object has a port open. The output buffer holds a maximum of 512 bytes.

Like the **write()** system function, **Write()** returns the actual number of bytes written, which will never be more than *numBytes*, and may be 0. If it fails (for example, if the BSerialPort object doesn't have a serial port open) or if it's interrupted before it can write anything, it returns **B_ERROR** (−1).

See also: **Read()**, **Open()**

Constants and Defined Types

This section lists the constants and types defined for the Device Kit, which currently contains only the BSerialPort class. Everything listed here is explained more fully in the descriptions of BSerialPort functions.

Constants

data_bits Constants

<device/SerialPort.h>

Enumerated constant

B_DATA_BITS_7

B_DATA_BITS_8

These constants name the possible number of data bits in a serial unit.

See also: **BSerialPort::SetDataBits()**

data_rate Constants

<device/SerialPort.h>

Enumerated constant

B_0_BPS

B_50_BPS

B_75_BPS

B_110_BPS

B_134_BPS

B_150_BPS

B_200_BPS

B_300_BPS

B_600_BPS

Enumerated constant

B_1200_BPS

B_1800_BPS

B_2400_BPS

B_4800_BPS

B_9600_BPS

B_19200_BPS

B_38400_BPS

B_57600_BPS

B_115200_BPS

These constants give the possible rates—in bits per second (bps)—at which data can be transmitted and received over a serial connection.

See also: **BSerialPort::SetDataRate()**

Flow Control Constants

<device/SerialPort.h>

Enumerated constant

B_SOFTWARE_CONTROL
B_HARDWARE_CONTROL

These constants form a mask that records the method(s) of flow control the serial port driver should use.

See also: **BSerialPort::SetFlowControl()**

parity_mode Constants

<device/SerialPort.h>

Enumerated constant

B_NO_PARITY
B_ODD_PARITY
B_EVEN_PARITY

These constants list the possibilities for parity when transmitting data over a serial connection.

See also: **BSerialPort::SetDataBits()**

stop_bits Constants

<device/SerialPort.h>

Enumerated constant

B_STOP_BITS_1
B_STOP_BITS_2

These constants name the possible number of stop bits in a serial unit.

See also: **BSerialPort::SetDataBits()**

Defined Types

data_bits

```
<device/SerialPort.h>
```

```
typedef enum { . . . } data_bits
```

This type is used to set and return the number of data bits in a serial unit.

See also: “**data_bits Constants**” above, `BSerialPort::SetDataBits()`

data_rate

```
<device/SerialPort.h>
```

```
typedef enum { . . . } data_rate
```

This type is used to set and return the rate at which data is sent and received through a serial connection.

See also: “**data_rate Constants**” above, `BSerialPort::SetDataRate()`

parity_mode

```
<device/SerialPort.h>
```

```
typedef enum { . . . } parity_mode
```

This type is used to set and return the type of parity that should be used when sending and receiving data.

See also: “**parity_mode Constants**” above, `BSerialPort::SetDataBits()`

stop_bits

```
<device/SerialPort.h>
```

```
typedef enum { . . . } stop_bits
```

This type is used to set and return the number of stop bits in a serial unit.

See also: “**stop_bits Constants**” above, `BSerialPort::SetDataBits()`

Developing a Device Driver

A device driver ties an input/output hardware device to the computer's operating system. To develop a driver, you have to know about both ends of that link:

- On the one hand, you need to be thoroughly familiar with the hardware device and its particular interface.
- On the other hand, you must understand the operating system and the demands it places on the driver.

Hardware specifications and manuals can provide you with the first kind of information; this documentation can help only with the second—that is, with information specific to the Be operating system. On the next page, you'll see a list of recommended documentation for the DMA controller, the PCI bus, and other hardware found inside the BeBox. This book is concerned solely with how a driver must be structured to work with Be system software.

Overview

On the BeBox, device drivers run as dynamically loaded add-on modules—as extensions of a host component of the operating system. The Application Server, the part of the operating system that's responsible for all graphics operations, is the host for graphics card drivers. The kernel acts as the host for all other drivers. The kernel and Application Server load drivers on demand, and can unload them when they're no longer needed.

Because drivers are linked to their hosts and run in the host's address space, they must play by the host's rules. The kernel and Application Server impose three different kinds of restrictions on loadable drivers:

- A driver must be constructed so that it can respond to its host. It has to be able to inform the host of the device or devices it drives, and it has to provide functions that the host can call to operate the driver. As an add-on module, a driver lacks an independent main thread of execution (and a `main()` function). Instead, it provides the host with entry points to driver functionality and responds to the host's instructions.
- A driver can call only those functions that it implements itself or that the host makes available to it. It cannot, for example, link to a library and call library functions. It links only to the kernel or Application Server and is limited to calling functions that the kernel or Application Server exports.
- The driver must be compiled as an add-on module and it must be located in a place where the host expects to find it.

Although the kernel and Application Server both impose these three kinds of restrictions on their loadable drivers, they each impose a different set of restrictions. The kernel's rules are not the Application Server's rules, and *vice versa*.

To learn the rules for developing a device driver that the kernel can load, begin with "Developing a Kernel-Loadable Driver" on page 19 next. To learn the Application Server's rules for graphics card drivers, begin with "Developing a Driver for a Graphics Card" on page 55.

Recommended Reading

For information on the PCI bus:

PCI Local Bus Specification, revision 2.1, June 1, 1995, PCI Special Interest Group, PO Box 14070, Portland OR 97214, (800) 433-5177 or (503) 797-4207

For information on the ISA bus, ISA 8259 interrupt controller, and ISA-standard 8237 DMA controller:

82378 System I/O (SIO), August 1994, order number 290473-004, Intel Corporation, 2200 Mission College Boulevard, PO Box 58119, Santa Clara, CA 95052

For information on the SCSI common access method (CAM):

Draft Proposed American National Standard *SCSI-2 Common Access Method Transport and SCSI Interface Module*, ASC X3T-10, revision 12, December 14, 1994, American National Standards Institute, 11 West 42nd Street, New York, NY 10036

Developing a Kernel-Loadable Driver

At the most basic level, devices (other than graphics cards) are controlled by system calls that the kernel traps and translates for the driver. Five different kinds of functions control input/output devices on the BeBox:

<code>open()</code>	Opens a device for reading or writing.
<code>close()</code>	Closes a device that was previously opened.
<code>read()</code>	Reads data from the device.
<code>write()</code>	Writes data to the device.
<code>ioctl()</code>	Formats, initializes, queries, and otherwise controls the device.

All these functions, except `ioctl()`, are Posix-compliant. Instead of `ioctl()`, Posix defines a set of functions like `tcsetattr()` and `tcflush()` to control data terminals. These functions are supported, but they can be treated as special cases of `ioctl()`. (Posix also defines a `fcntl()` function for file control that has the same syntax as `ioctl()`.)

When `open()`, `close()`, `read()`, `write()`, or `ioctl()` is called for a device, the kernel expects a driver to do the work that's required. Each driver must implement a set of functions that correspond directly to the five system calls. Everything the driver does to operate the device is initiated through one of these functions.

Because drivers run in the kernel's address space as extensions of the kernel, they must conform to the kernel's expectations. The sections below discuss the three types of restrictions that the kernel imposes on drivers:

- “Entry Points” below describes how drivers must be constructed to work with the kernel. The driver must provide the kernel with entry points to its functionality and follow the kernel's instructions.
- “Exported Functions” on page 27 discusses the kinds of functions that the kernel exports for drivers. These include some from the Kernel Kit, Support Kit, and standard C library, and some that are defined especially for drivers. The driver is limited to calling functions that it itself implements or that the kernel exports.
- “Installation” on page 32 discusses how to compile a driver and install it on the BeBox. The driver must be installed where the kernel can find it.

Specific documentation on the exported functions that the kernel defines just for drivers follows these sections.

Entry Points

The kernel loads a device driver when it's needed—typically when someone first attempts to open the device for reading or writing. Opening a device is a prerequisite to using it.

To the `open()` function, drivers are identified by a fictitious pathname beginning with `/dev/`. For example, this code opens the parallel port driver for writing:

```
int fd = open("/dev/parallel", O_WRONLY);
```

The first thing the kernel must do is match the device name—“`/dev/parallel`” in this case—to a driver; it must find a driver for the device. The driver might be one that has already been loaded, or it might be one that the kernel must search for and load. Loadable drivers reside in the `/system/drivers` directory; this is where the kernel looks for drivers and where they all must be installed.

Once a driver has been located and loaded, the kernel begins communicating with it—first to get information from it and test whether it's the right driver, then to initialize it and have it open the device.

One key piece of information that the kernel needs from the driver is the names of all its devices. Another is a list of the functions it can call to exercise those devices. For each device, the driver implements a set of hook functions that open the device, control it, read data from it, write data to it, and perhaps eventually close it. These hooks correspond to the system functions discussed above.

To give the kernel initial access to this information, drivers make declarations—of data structures and functions—using names the kernel will look for. Four such names will be discussed in the following sections:

<code>init_driver()</code>	Initializes the driver when it's loaded.
<code>devices</code>	Declares the devices and their hook functions.
<code>publish_device_names()</code>	Lists the devices the driver handles.
<code>find_device_entry()</code>	Associates a device with its hook functions.

These are the main entry points for driver control.

Driver Initialization

Immediately after loading a driver, the kernel gives it a chance to initialize itself. If the driver implements a function called `init_driver()`, the kernel will call it before proceeding with anything else—before asking the driver to open a device. The function should expect no arguments and return either `B_ERROR` or `B_NO_ERROR`:

```
long init_driver(void)
```

A return of `B_ERROR` means that the driver can't continue; the kernel will consequently unload it. A return of `B_NO_ERROR` means that all is well; the kernel will continue by

asking the driver to open a device. The absence of an `init_driver()` function is equivalent to a return of `B_NO_ERROR`.

`init_driver()` might go a long way toward initializing the data structures the driver uses to do its work. However, details specific to a particular device should be left to the hook function that opens that device.

Driver initialization happens just once—when the driver is loaded. In contrast, devices might be opened and closed many times while the driver continues to reside in the kernel.

Device Declarations

For the kernel to find the driver for a given device, all drivers must declare the names of the devices they control. For the kernel to be able to communicate with the driver to operate the device, every driver must declare a set of device-specific hook functions the kernel can call.

These declarations are made in a `device_entry` structure that maps the device name to the set of hook functions. This structure is declared in **device/Drivers.h** and contains the following fields:

<code>const char *name</code>	The name of the device—for example, “/dev/serial”. This is the same name that’s passed to <code>open()</code> . Driver code can assign any name it wants to the device, but it must begin with the “/dev/” prefix, which distinguishes devices from ordinary files.
<code>device_open_hook open</code>	The function that the kernel should call to open the device. It corresponds to the <code>open()</code> system call.
<code>device_close_hook close</code>	The function that should be called to close the device. It corresponds to the <code>close()</code> system call.
<code>device_control_hook control</code>	The function that the kernel should call to control the device, including querying the driver for information about it. The kernel will invoke this function to respond to <code>ioctl()</code> calls.
<code>device_io_hook read</code>	The function that should be called to read data from the device. The kernel will invoke this function to respond to the <code>read()</code> system call.
<code>device_io_hook write</code>	The function that should be called to write data to the device. The kernel will invoke this function to respond to <code>write()</code> system calls.

(The five functions are described in more detail under “Hook Functions” below.)

A driver declares one **device_entry** structure for each device it can drive. If it can handle more than one device, it must provide a **device_entry** structure for each one. If it permits a device to be referred to by more than one name, it must provide a structure for each name it recognizes.

There are two ways for a driver to provide the kernel with the information in a **device_entry** structure: If the list of devices is known at compile time, the driver can declare them statically. If the list might change at run time, it can return them dynamically.

Static Drivers

Most drivers are designed to handle a fixed set of known devices—perhaps a single device or perhaps many. Such drivers should declare a null-terminated array of **device_entry** structures under the global name **devices**:

```
device_entry devices[]
```

For example, the serial port driver might declare a **devices** array that looks like this:

```
device_entry devices[7] = {
    {"/dev/serial1", open_func, close_func, control_func,
      read_func, write_func},
    {"/dev/serial2", open_func, close_func, control_func,
      read_func, write_func},
    {"/dev/serial3", open_func, close_func, control_func,
      read_func, write_func},
    {"/dev/serial4", open_func, close_func, control_func,
      read_func, write_func},
    {"/dev/com3", open_com_func, close_func, control_com_func,
      read_func, write_func},
    {"/dev/com4", open_com_func, close_func, control_com_func,
      read_func, write_func},
    0
};
```

In this case, the driver handles the four serial ports seen on the back of the BeBox, each of which it identifies by a different name. It can also control “com3” and “com4” ports on an add-on board.

As this example illustrates, the hook functions declared in a **device_entry** structure are specific to the device. For the most part, the serial port driver above uses the same set of functions to operate all the devices, but declares special functions for opening and controlling “com3” and “com4”.

Note also that the array is null terminated.

Dynamic Drivers

A driver can also provide **device_entry** information dynamically. Instead of a **devices** array, it implements two functions, **publish_device_names()** and **find_device_entry()**.

The first of these functions should be declared as follows:

```
char **publish_device_names(const char *deviceName)
```

If passed a proposed *deviceName* that matches the name of a device the driver handles, or if passed a **NULL** device name, this function should return a null-terminated array of all the names of all the devices that it handles. For example, the serial port driver described above would return the following array:

```
"/dev/serial1",
"/dev/serial2",
"/dev/serial3",
"/dev/serial4",
"/dev/com3",
"/dev/com4",
0
```

However, if the proposed device name doesn't match any that the driver handles, **publish_device_names()** should return **NULL**.

While **publish_device_names()** informs the kernel of the devices that the driver handles, **find_device_entry()** returns entry information about a particular device. It has the following form:

```
device entry *find_device_entry(const char *deviceName)
```

If passed the name of a device that the driver knows about, this function should return the **device_entry** for that name. If the *deviceName* doesn't match one of the driver's devices, it should return **NULL**.

The kernel first calls **publish_device_names()** during the boot sequence to find what devices the driver handles. It may call the function again to update the list when it tries to match a driver to a specific device. If a match is made, it calls **find_device_entry()** to get the list of hook functions for the device.

Hook Functions

The five hook functions that are declared in a **device_entry** structure can have any names that you want to give them, provided that they don't clash with names that the kernel exports (see "Exported Functions" on page 27). However, their syntax is strictly prescribed by the kernel (through type definitions found in **device/Drivers.h**).

The five functions have two points in common: First, they each return an error code, which should be 0 (**B_NO_ERROR**) if there is no error. The error value is passed through as the return value for the **open()**, **close()**, **read()**, **write()**, or **ioctl()** system call that caused

the kernel to invoke the driver function. The driver should return error values that are compatible with ones that are expected from those functions.

Second, all five functions are passed information identifying the device. As their first argument, they receive a pointer to a **device_info** structure (also defined in **device/Drivers.h**), which contains just two fields:

<code>device_entry *entry</code>	The device_entry structure for the device that's being operated on. This is a copy of information that the driver declared in its devices array or that its find_device_entry() function returned.
<code>void *private_data</code>	Arbitrary data that describes the device. This data is a way for the driver to record information about the device and have it persist between function calls. Although the kernel stores this data and passes it to the driver, the driver initializes it and maintains it; the kernel doesn't query or modify it.

Thus, all of the device-specific hook functions have the same return types and initial arguments:

```
long function(device_info *info, . . .)
```

Differences among the functions are discussed below.

Opening and Closing a Device

The function that opens a device is of type **device_open_hook** and the one that closes it is of type **device_close_hook**. They're defined as follows:

```
typedef long (*device_open_hook)(device_info *info, ulong flags)
typedef long (*device_close_hook)(device_info *info)
```

The *flags* mask that is passed to the **open()** system call is passed through to the device function. It typically will contain a flag like **O_RDONLY**, **O_RDWR**, or **O_WRONLY**.

Since the hook function that opens the device is the first one that's called, it might set up the **device_info** description of the device (to the extent that **init_driver()** hasn't already done so). It might also use that description, or some static data, to record whether or not the device is currently open. Typically, only one process can have a device open at a time. If the hook function sees that the device is already open, it can refuse to open it again.

Whatever values these functions return will also be returned by the **open()** and **close()** system calls.

Reading and Writing Data

The driver functions that read data from and write data to a device must be of type **device_io_hook**, which is defined as follows:

```
typedef long (*device_io_hook)(device_info *info,
                               void *buffer, ulong numBytes, ulong position)
```

The function that reads data from the device should place up to *numBytes* of data into the specified *buffer* beginning at *position*. The hook that writes to the device should take *numBytes* of data from the *buffer* beginning at *position*.

Whatever values these functions return will also be returned by the **read()** and **write()** system calls.

Controlling the Device

The hook function that initializes, formats, queries, and otherwise controls a device is of type **device_control_hook**, defined as follows:

```
typedef long (*device_control_hook)(device_info *info, ulong op, void *data)
```

The second argument, *op*, is a constant that specifies the particular control operation that the function should perform. The third argument, *data*, points either to some information that the control function needs to carry out the *op* operation or to a data structure that it should fill in with information that the operation requests. The interpretation of the data pointer depends entirely on the nature of the operation and will differ from operation to operation; the *op* and *data* arguments go hand-in-hand.

For example, if the *op* code is **SET_CONFIG**, *data* might point to a structure with values that the control function should use to re-configure the device. If the operation is **GET_ENABLED_STATE**, *data* might point to an integer that the function would be expected to set to either 1 or 0. If it's **RESTART**, *data* might simply be **NULL**.

The kernel defines a number of control operations (which are explained in the next section). These are operations that the kernel might call upon any driver to perform.

If you define your own control operations (for an **ioctl()** call on your driver), you should be sure that they aren't confused with any that the kernel currently defines—or any that it will define in the future. We pledge that all system-defined control constants will have values below **B_DEVICE_OP_CODES_END**. The constants you define should be increments above this value. For example:

```
enum {
    REPORT_STATUS = B_DEVICE_OP_CODES_END + 1,
    SET_TIMER,
    . . .
}
```

If a control function doesn't recognize the *op* code it's passed or can't perform the requested operation, it should return **B_ERROR** (−1).

Control Operations

Several control operations are defined by the kernel. The kernel can request any driver to perform these operations, even in the absence of an **ioctl()** call. A control function should respond to as many of these requests as it can. It should respond to inappropriate or unrecognized requests by returning **B_ERROR**.

The set of system-defined control operations is described below.

B_GET_SIZE and **B_SET_SIZE**

These control operations request the driver to get and set the memory capacity of the physical device. The capacity is measured in bytes and is recorded as a **ulong** integer. For a **B_GET_SIZE** request, the control function should write this number to the location referred to by the *data* pointer. For **B_SET_SIZE**, *data* will be the requested number of bytes (not a pointer to it).

B_SET_BLOCKING_IO and **B_SET_NONBLOCKING_IO**

These operations determine whether or not the driver should block when reading and writing data. **B_SET_BLOCKING_IO** requests the driver to put itself in blocking mode. Its read function should wait for data to arrive if none is readily available and its write function should wait for the device to be ready to accept data if it's not immediately free to take it. If **B_SET_NONBLOCKING_IO** is requested, the read function should return immediately if there is no data available to read and the write function should return immediately if the device isn't ready to accept written data.

For these operations, the *data* argument doesn't contain a meaningful value.

B_GET_READ_STATUS and **B_GET_WRITE_STATUS**

These control operations request the driver to report whether or not it's ready to read and write without blocking. The control function should respond by placing **TRUE** or **FALSE** as a **ulong** integer in the location that *data* points to.

For **B_GET_READ_STATUS**, it should respond **TRUE** if there's data waiting to be read, and **FALSE** if not. For **B_GET_WRITE_STATUS**, it should respond **TRUE** if the device is free to accept data, and **FALSE** if not.

B_GET_GEOMETRY

This *op* code requests the driver to supply information about the physical configuration of the device; it's generally appropriate only for mass storage devices. The control function should write the requested information into the **device_geometry** structure that the *data* pointer refers to. A **device_geometry** structure contains the following fields:

ulong bytes_per_sector	The number of bytes in each sector of storage.
ulong sectors_per_track	The number of sectors in each track.
ulong cylinder_count	The number of cylinders.
ulong head_count	The number of heads.
bool removable	Whether or not the storage medium can be removed (TRUE if it can be, FALSE if not).
bool read_only	Whether or not the medium can be read but not written (TRUE if it cannot be written, FALSE if it can).
bool write_once	Whether or not the medium can be written once, after which it becomes read-only (TRUE if it can be written only once, FALSE if it cannot be written or can be written more than once).

B_FORMAT

This operation requests the control function to format the device. The *data* argument doesn't contain any valid information.

Exported Functions

After a driver has been loaded, it runs as part of the kernel in the kernel's address space. It therefore is restricted to calling functions (a) that it implements or (b) that the kernel makes available to it. The driver links to the kernel alone; it cannot also independently link to something else, even the standard C library.

The kernel exports five kinds of functions so that they're available to a driver:

- Support Kit functions, such as **read_32_swap()** and **atomic_or()**. Like the error constants and data types that are defined in the Support Kit, these functions are available to drivers.
- Standard kernel functions, such as **area_for()** and **write_port()**, that were documented in the chapter on the Kernel Kit. Because the driver runs in the kernel's address space, it accesses these functions directly, not through the Kit. For this

reason, not all of the functions are available to drivers; there are some services the kernel can provide to others, but not to itself.

- Library functions that the kernel incorporates. These are functions from the standard C library that have been adopted by the kernel and that the kernel, in turn, exports to the driver. They're a small, selected subset of library functions.
- Standard system calls, such as `read()` and `ioctl()`.
- Special functions that are implemented specifically for device drivers.

Functions from all five groups are listed in the sections below. Special driver functions are documented in detail in the section entitled "Functions for Drivers" on page 33.

Support Kit Functions

The kernel exports the following Support Kit functions:

<code>read_16_swap()</code>	<code>atomic_and()</code>
<code>read_32_swap()</code>	<code>atomic_or()</code>
<code>write_16_swap()</code>	<code>atomic_add()</code>
<code>write_32_swap()</code>	<code>real_time_clock()</code>

See the chapter on the Support Kit for descriptions of these functions.

Kernel Kit Functions

Most functions from the Kernel Kit are available to drivers. However, a few are not, sometimes because it would make no sense for a driver to call the function, and sometimes because it's difficult for the kernel to provide its very basic services to its own modules. In some cases, a special function is defined for drivers that takes the place of the missing Kit function. For example, `spawn_thread()` can't spawn a thread in the kernel. Since drivers run in the kernel, they need to use the special `spawn_kernel_thread()` instead. Similarly, `debugger()` can't be used to debug the kernel. Drivers should call `kernel_debugger()` instead.

The following Kernel Kit functions are exported for drivers:

<u>Semaphores</u>	
<code>create_sem()</code>	<code>get_sem_count()</code>
<code>delete_sem()</code>	<code>set_sem_owner()</code>
<code>get_sem_info()</code>	<code>release_sem()</code>
<code>get_nth_sem_info()</code>	<code>release_sem_count()</code>

Threads

find_thread()	suspend_thread()
rename_thread()	resume_thread()
set_thread_priority()	wait_for_thread()
get_thread_info()	exit_thread()
get_nth_thread_info()	kill_thread()

Teams

kill_team()
get_team_info()
get_nth_team_info()

Ports

create_port()	read_port()
delete_port()	write_port()
find_port()	port_count()
get_port_info()	port_buffer_size()
get_nth_port_info()	set_port_owner()

Time

snooze()
system_time()

Other

area_for()
get_system_info()

C Library Functions

The kernel exports a small number of functions from the standard C library. They include:

Functions declared in stdlib.h

atof()	malloc()
atoi()	calloc()
atol()	free()
strtod()	abs()
strtol()	div()
strtoul()	labs()
bsearch()	ldiv()
qsort()	

Functions and macros declared in ctype.h

isalnum()	ispunct()
isalpha()	isspace()
iscntrl()	isupper()
isdigit()	isxdigit()
isgraph()	
islower()	tolower()
isprint()	toupper()

Functions declared in string.h

strcspn()	memchr()
strlen()	memcmp()
strncat()	memcpy()
strncmp()	memmove()
strncpy()	memset()
strpbrk()	
strchr()	
strspn()	
strstr()	

Functions declared in stdio.h

```
sprintf()
vsprintf()
```

The driver accesses these functions from the kernel, not from the library.

System Calls

The kernel also exports the five system calls that control devices:

```
open()
close()
read()
write()
ioctl()
```

Kernel Functions for Drivers

The kernel defines the following functions especially for drivers. For full documentation of these functions, see “Functions for Drivers” on page 33.

Semaphores:

acquire_sem_flag()	acquire_sem_count_timeout_flag()
acquire_sem_count_flag()	release_sem_count_flag()
acquire_sem_timeout_flag()	

Spinlocks:

acquire_spinlock()
release_spinlock()

Disabling interrupts:

disable_interrupts()
restore_interrupts()

Interrupt handling:

set_io_interrupt_handler()
disable_io_interrupt()
enable_io_interrupt()

set_isa_interrupt_handler()
disable_isa_interrupt()
enable_isa_interrupt()

Memory management:

lock_memory()
unlock_memory()

get_memory_map()
isa_address()

ISA DMA:

start_isa_dma()
start_scattered_isa_dma()
make_isa_dma_table()

lock_isa_dma_channel()
unlock_isa_dma_channel()

PCI:

read_pci_config()
write_pci_config()
get_nth_pci_info()

Debugging:

dprintf()
set_dprintf_enabled()
kernel_debugger()

Hardware versions:

motherboard_version()
io_card_version()

SCSI common access method:

xpt_init()
xpt_ccb_alloc()
xpt_ccb_free()

xpt_action()
xpt_bus_register()
xpt_bus_deregister()

Other

spin()
spawn_kernel_thread()

Installation

The driver must be compiled as an add-on image, as explained in the Kernel Kit chapter. For the kernel to be able to find the compiled driver, it must be installed in the **/system/drivers** directory. This is the only place that the kernel looks for drivers to load.

When a device is opened, the kernel first looks for its driver among those that are already loaded. Failing that, it looks on a floppy disk (in **/fd/system/drivers**). Failing to find one there, it looks next on the boot disk (in **/boot/system/drivers**).

If the **/system/drivers** directory contains more than one driver for the same device, it's indeterminate which one will be loaded.

You can give your driver any name you wish.

Functions for Drivers

The kernel exports a number of functions for the benefit of device drivers. These are functions that drivers can call to do their work; they're not functions that are available to applications. Although implemented by the kernel, they're not part of the Kernel Kit. The device driver accesses these functions directly from the kernel, not through a library.

`acquire_sem_flag()`, `acquire_sem_count_flag()`,
`acquire_sem_timeout_flag()`, `acquire_sem_count_timeout_flag()`,
`release_sem_count_flag()`

`<device/KernelExport.h>`

`long acquire_sem_flag(sem_id sem, long flags)`

`long acquire_sem_count_flag(sem_id sem, long count, long flags)`

`long acquire_sem_timeout_flag(sem_id sem, double timeout, long flags)`

`long acquire_sem_count_timeout_flag(sem_id sem, long count, double timeout,
long flags)`

`long release_sem_count_flag(sem_id sem, long count, long flags)`

These functions are equivalent to their Kernel Kit counterparts, except that their names end in “_flag” and they take an additional *flags* argument. The *flags* argument is a mask that can be 0 or any combination of three constants:

`B_CAN_INTERRUPT`

`B_DO_NOT_RESCHEDULE`

`B_CHECK_PERMISSION`

The `B_CAN_INTERRUPT` flag is appropriate for the various `acquire_sem. . .flag()` functions; it indicates that the function can be interrupted by a signal. If the signal arrives while the function is blocked waiting for the semaphore, it returns `B_INTERRUPTED` immediately. If the signal is pending before the function call and the semaphore is available, it acquires the semaphore and returns `B_NO_ERROR`. However, if the semaphore isn't available, it returns `B_INTERRUPTED` immediately; it won't block.

`B_CAN_INTERRUPT` is the normal behavior for the Kernel Kit's `acquire_sem. . .()` functions, but drivers must ask for it explicitly. If it isn't included in the *flags* mask, the `acquire_sem. . .flag()` functions will block while waiting for a semaphore even in the presence of a signal.

You may want to avoid using the `B_CAN_INTERRUPT` flag if the semaphore protects a critical section of code that can't easily be undone in case it's interrupted. However, if you

do use the flag and the semaphore is never released, your thread will block forever waiting to acquire it.

The Kernel Kit's `acquire_sem. . .()` functions are not exported for drivers; drivers should use the `acquire_sem. . .flag()` functions to acquire semaphores instead.

`B_DO_NOT_RESCHEDULE` is an appropriate flag for `release_sem_count_flag()`. When one thread releases a semaphore, another thread that has been blocked waiting to acquire the semaphore becomes free to continue its activity. Normally, the scheduler facilitates this process by automatically rescheduling threads when one of them releases a semaphore that another thread is waiting for. `B_DO_NOT_RESCHEDULE` prevents this so that, for example, a single thread can release a number of semaphores before any automatic rescheduling takes place.

It's strongly recommended that interrupt handlers always include `B_DO_NOT_RESCHEDULE` in the *flags* mask when they release a semaphore. Rescheduling while an interrupt is being handled is generally not a good idea.

In addition to `release_sem_count_flag()`, two Kernel Kit functions, `release_sem()` and `release_sem_count()`, are exported for drivers. These functions don't prevent automatic rescheduling.

`B_CHECK_PERMISSION` is appropriate for both the `acquire_sem. . .flag()` functions and `release_sem_count_flag()`. Drivers, because they run in the kernel, have access to all semaphores, even those that the kernel uses for its own purposes. This flag places the system semaphores off-limits; they won't be acquired or released.

See also: `acquire_sem()` in the Kernel Kit

`acquire_spinlock()`, `release_spinlock()`

<device/KernelExport.h>

`void acquire_spinlock(spinlock *lock)`

`void release_spinlock(spinlock *lock)`

These functions acquire and release the *lock* spinlock. Spinlocks, like semaphores, are used to protect critical sections of code that must remain on the same processor for a single path of execution—for example, code that atomically accesses a hardware register or a shared data structure. A common use for spinlocks is to protect data structure that both an interrupt handler and normal driver code must access.

However, spinlocks work quite differently from semaphores. No count is kept of how many times a thread has acquired the lock, for example, so calls to `acquire_spinlock()` and `release_spinlock()` should not be nested. More importantly, `acquire_spinlock()` spins while attempting to acquire the lock; it doesn't block or release its hold on the CPU.

These functions assume that interrupts have been disabled. They should be nested within calls to `disable_interrupts()` and `restore_interrupts()` as follows:

```
spinlock lock;
cpu_status former = disable_interrupts();
acquire_spinlock(&lock);
/* critical code goes here */
release_spinlock(&lock);
restore_interrupts(former);
```

These two pairs of functions enable the thread to get into the critical code without rescheduling. Disabling interrupts ensures that the thread won't be preemptively rescheduled. Because `acquire_spinlock()` doesn't block, it also won't be voluntarily rescheduled.

Executing the critical code under the protection of the spinlock guarantees that no other thread will execute the same code at the same time on another processor. Spinlocks should be held only as long as necessary and released as quickly as possible.

See also: `create_spinlock()`

`create_spinlock()`, `delete_spinlock()`

```
<device/KernelExport.h>
spinlock *create_spinlock(void)
void delete_spinlock(spinlock *lock)
```

< These functions will, when implemented and exported, produce and destroy spinlocks. Currently, they're declared but not exported. To create a spinlock at present, simply declare a `spinlock` variable and pass a pointer to it to `acquire_spinlock()`. >

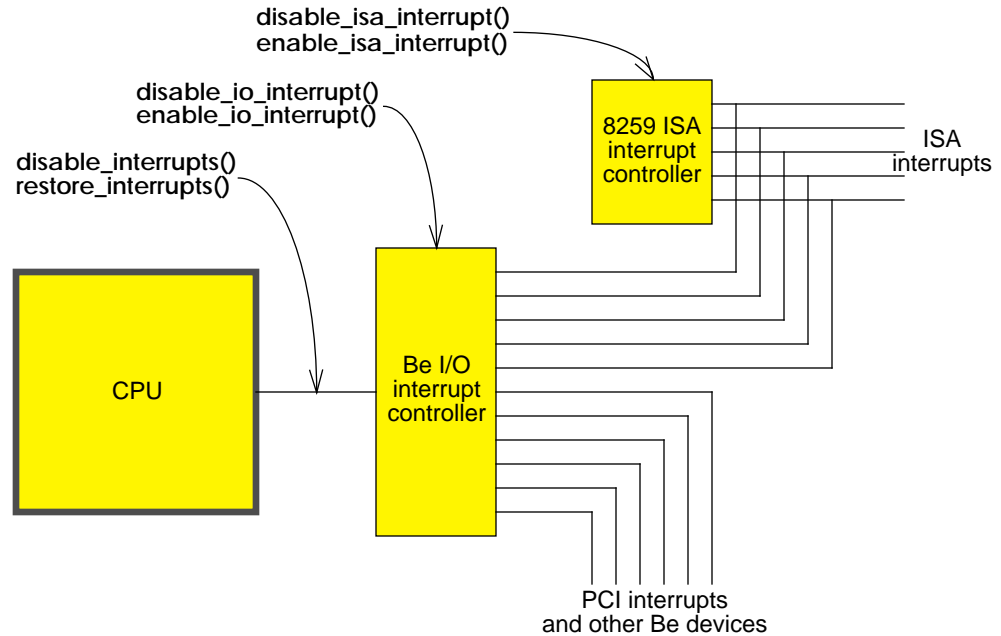
See also: `acquire_spinlock()`

`disable_interrupts()`, `restore_interrupts()`

```
<device/KernelExport.h>
cpu_status disable_interrupts(void)
void restore_interrupts(cpu_status status)
```

These functions disable interrupts at the CPU (the one the caller is currently running on) and restore them again. `disable_interrupts()` prevents the CPU from being interrupted and returns its previous status—whether or not interrupts were already disabled before the `disable_interrupts()` call. `restore_interrupts()` restores the previous *status* of the CPU, which should be the value that `disable_interrupts()` returned. Passing the status returned by `disable_interrupts()` to `restore_interrupts()` allows these functions to be paired and nested.

As diagrammed below, individual interrupts can be enabled and disabled at two other hardware locations. `disable_isa_interrupt()` and `enable_isa_interrupt()` work at the ISA standard 8259 interrupt controller, and `disable_io_interrupt()` and `enable_io_interrupt()` act at the Be-defined I/O interrupt controller that combines ISA and PCI interrupts.



Interrupts that have been disabled by `disable_interrupts()` must be reenabled by `restore_interrupts()`.

See also: `acquire_spinlock()`, `set_io_interrupt_handler()`, `set_isa_interrupt_handler()`

`disable_io_interrupt()` see `set_io_interrupt_handler()`

`disable_isa_interrupt()` see `set_isa_interrupt_handler()`

`dprintf()`, `set_dprintf_enabled()`, `kernel_debugger()`

<device/KernelExport.h>

void `dprintf`(const char *format, ...)

bool `set_dprintf_enabled`(bool isEnabled)

void `kernel_debugger`(const char *string)

`dprintf()` is a debugging function that has the same syntax and behavior as standard C `printf()`, except that it writes its output to the fourth serial port (“/dev/serial4”) at a data rate of 19,200 bits per second. By default, `dprintf()` is disabled.

`set_dprintf_enabled()` enables `dprintf()` if the *isEnabled* flag is `TRUE`, and disables it if the flag is `FALSE`. It returns the previous enabled state. Calls to this function can be nested by caching the return value of a call that disables printing and passing it to the paired call that restores the previous state.

`kernel_debugger()` drops the calling thread into a debugger that writes its output to the fourth serial port at 19,200 bits per second, just as `dprintf()` does. This debugger produces *string* as its first message; it's not affected by `set_dprintf_enabled()`.

`kernel_debugger()` is identical to the `debugger()` function documented in the Kernel Kit, except that it works in the kernel and engages a different debugger. Drivers should use it instead of `debugger()`.

See also: `debugger()` in the Kernel Kit

`enable_io_interrupt()` see `set_io_interrupt_handler()`

`enable_isa_interrupt()` see `set_isa_interrupt_handler()`

`get_memory_map()`

<device/KernelExport.h>

```
long get_memory_map(void *address, ulong numBytes, physical_entry *table,
                    long numEntries)
```

Locates the separate pieces of physical memory that correspond to the contiguous buffer of virtual memory beginning at *address* and extending for *numBytes*. Each piece of physical memory is described by a `physical_entry` structure. It has just two fields:

<code>void *address</code>	The address of a block of physical memory.
<code>ulong size</code>	The number of bytes in the block.

This function is passed a pointer to a *table* of `physical_entry` structures. It fills in the table, stopping when the entire buffer of virtual memory has been described or when *numEntry* entries in the table have been written, whichever comes first.

If the *table* provided isn't big enough, you'll need to call `get_memory_map()` again and ask it to describe the rest of the buffer. If the table is too big, this function sets *size* field of the entry following the last one it needed to 0. This indicates that it has finished mapping the entire *address* buffer.

Memory should be locked while it is being mapped. Before calling `get_memory_map()`, call `lock_memory()` to make sure that it all stays in place:

```
physical_entry table[count];
lock_memory(someAddress, someNumberOfBytes, FALSE);
get_memory_map(someAddress, someNumberOfBytes, table, count);
. . .
unlock_memory(someAddress, someNumberOfBytes);
```

< This function consistently returns `B_NO_ERROR`. >

See also: `lock_memory()`, `start_isa_dma()`

`get_nth_pci_info()`

<device/PCI.h>

```
long get_nth_pci_info(long index, pci_info *info)
```

This function looks up the PCI device at *index* and provides a description of it in the `pci_info` structure that *info* refers to. Indices begin at 0 and there are no gaps in the list.

The `pci_info` structure contains a number of fields that report values found in the configuration register space for the device and it also describes how the device has been mapped into the system. The following fields are common to all devices:

<code>ushort vendor_id</code>	An identifier for the manufacturer of the device.
<code>ushort device_id</code>	An identifier for the particular device of the vendor, assigned by the vendor.
<code>uchar bus</code>	The bus number.
<code>uchar device</code>	The number that identifies the location of the device on the bus.
<code>uchar function</code>	The function number in the device.
<code>uchar revision</code>	A device-specific version number, assigned by the vendor.
<code>uchar class_api</code>	The type of specific register-level programming interface for the device (the lower byte of the class code field).
<code>uchar class_sub</code>	The specific type of function the device performs (the middle byte of the class code field).
<code>uchar class_base</code>	The broadly-defined device type (the upper byte of the class code field).

<code>uchar line_size</code>	The size of the system cache line, in units of 32-bit words.
<code>uchar latency</code>	The latency timer for the PCI bus master.
<code>uchar header_type</code>	The header type.
<code>uchar bist</code>	The contents of the register for the built-in self test.
<code>uchar u</code>	A union of structures, one for each header type.

Currently, there's only one header (type 0x00), but in the future there may be others. Consequently, header-specific information is recorded in a union of structures, one for each header type. The union (named simply `u`) at present has just one member, a structure for the current header (named `h0`):

```
typedef struct {
    ...
    union {
        struct {
            ...
        } h0;
    } u;
} pci_info
```

The fields of the `h0` structure are:

<code>ulong cardbus_cis</code>	The CardBus CIS pointer.
<code>ushort subsystem_id</code>	The vendor-assigned identifier for the add-in card containing the device.
<code>ushort subsystem_vendor_id</code>	The identifier for manufacturer of the add-in card that contains the device.
<code>ulong rom_base</code>	The base address for the expansion ROM, as viewed from the host processor.
<code>ulong rom_base_pci</code>	The base address for the expansion ROM, as viewed from PCI.
<code>ulong rom_size</code>	The amount of memory in the expansion ROM, in bytes.
<code>ulong base_registers[6]</code>	The base addresses of requested memory spaces and I/O spaces, as viewed from the host processor.
<code>ulong base_registers_pci[6]</code>	The base addresses of requested memory spaces and I/O spaces, as viewed from PCI.
<code>ulong base_register_sizes[6]</code>	The sizes of requested memory spaces and I/O spaces.

uchar base_register_flags [6]	The flags from the base-address registers.
uchar interrupt_line	The interrupt line. This number identifies the interrupt associated with the device. See <code>set_io_interrupt_handler()</code> .
uchar interrupt_pin	The interrupt pin that the device uses.
uchar min_grant	The minimum burst period the device needs, assuming a clock rate of 33 MHz.
uchar max_latency	The maximum frequency at which the device needs access to the PCI bus.

In **device/PCI.h**, you'll find a number of constants that you can use to test various fields of a `pci_info` structure. See the *PCI Local Bus Specification*, published by the PCI Special Interest Group (Portland, OR) for more information on the configuration of a PCI device.

`get_nth_pci_info()` returns `B_NO_ERROR` if it successfully describes a PCI device, and `B_ERROR` if it can't find the device (for example, if *index* is out-of-range).

See also: `read_pci_config()`

io_card_version() see **motherboard_version()**

isa_address()

<device/KernelExport.h>

void *isa_address(long *offset*)

Returns the virtual address corresponding to the specified *offset* in the ISA I/O address space. By passing an *offset* of 0, you can find the base address that's mapped to the ISA address space.

kernel_debugger() see **dprintf()**

lock_isa_dma_channel(), unlock_isa_dma_channel()

<device/KernelExport.h>

long lock_isa_dma_channel(long *channel*)

long unlock_isa_dma_channel(long *channel*)

These functions reserve an ISA DMA *channel* and release a channel previously reserved. They return `B_NO_ERROR` if successful, and `B_ERROR` if not. Like semaphores, these functions work only if all participating parties adhere to the protocol.

There are 7 ISA DMA channels. In general, they're used as follows:

<u>Channel</u>	<u>Use</u>
0	Unreserved, available
1	Unreserved, available
2	Reserved for the floppy disk controller
3	Reserved for the parallel port driver
5	Reserved for IDE
6	Reserved for sound
7	Reserved for sound

Channel 4 is taken by the system; it cannot be used.

lock_memory(), unlock_memory()

<device/KernelExport.h>

long **lock_memory**(void *address, ulong numBytes, bool willChange)

long **unlock_memory**(void *address, ulong numBytes)

lock_memory() makes sure that all the memory beginning at the specified virtual *address* and extending for *numBytes* is resident in RAM, and locks it so that it won't be paged out until **unlock_memory()** is called. It pages in any of the memory that isn't resident at the time it's called.

The *willChange* flag should be **TRUE** if any part of the memory range will be altered while it's locked—especially if the hardware device will do anything to modify the memory, since that won't otherwise be noticed by the system and the modified pages may not be written. The *willChange* flag should be **FALSE** if the memory won't change while it's locked.

Each of these functions returns **B_NO_ERROR** if successful and **B_ERROR** if not. The main reason that **lock_memory()** would fail is that you're attempting to lock more memory than can be paged in.

make_isa_dma_table() see start_isa_dma()

motherboard_version(), io_card_version()

<device/KernelExport.h>

long **motherboard_version**(void)

long **io_card_version**(void)

These functions return the current versions of the motherboard and of the I/O card.

read_pci_config(), write_pci_config()

<device/PCI.h>

```
long read_pci_config(uchar bus, uchar device, uchar function,
                    long offset, long size)
```

```
void write_pci_config(uchar bus, uchar device, uchar function,
                     long offset, long size, long value)
```

These functions read from and write to the PCI configuration register space. The *bus*, *device*, and *function* arguments can be read from the *bus*, *device*, and *function* fields of the *pci_info* structure provided by *get_nth_pci_info()*. They identify the configuration space that belongs to the device.

The *offset* is an offset to the location in the 256-byte configuration space that is to be read or written and *size* is the number of bytes to be read from that location or written to it. Permitted sizes are 1, 2, and 4 bytes. *read_pci_config()* returns the bytes that are read, *write_pci_config()* writes *size* bytes of *value* to the *offset* location.

See also: *get_nth_pci_info()*

release_sem_count_flag() see *acquire_sem_flag()*

release_spinlock() see *acquire_spinlock()*

restore_interrupts() see *disable_interrupts()*

set_dprintf_enabled() see *dprintf()*

**set_io_interrupt_handler(),
disable_io_interrupt(), enable_io_interrupt()**

<device/KernelExport.h>

```
long set_io_interrupt_handler(long interrupt,
                             interrupt_handler function, void *data)
```

```
long disable_io_interrupt(long interrupt)
```

```
long enable_io_interrupt(long interrupt)
```

These functions manage interrupts at the Be-designed I/O interrupt controller that combines ISA and PCI interrupts. The *interrupt* can be an ISA IRQ value or the *interrupt_line* field read from the *pci_info* structure provided by *get_nth_pci_info()*.

set_io_interrupt_handler() installs the handler *function* that will be called each time the specified *interrupt* occurs. This function should have the following syntax:

```
bool handler(void *data)
```

The *data* that's passed to `set_io_interrupt_handler()` will be passed to the handler function each time it's called. It can be anything that might be of use to the handler, or `NULL`. This function should always return `TRUE`.

`set_io_interrupt_handler()` itself returns `B_NO_ERROR` if successful in installing the handler, and `B_ERROR` if not.

`disable_io_interrupt()` disables the named *interrupt*, and `enable_io_interrupt()` reenables it. Both functions return `B_ERROR` for an invalid *interrupt* number, and `B_NO_ERROR` otherwise. Neither function takes into account the disabled or enabled state of the interrupt as it might be affected by other functions, such as `disable_isa_interrupt()` or `restore_interrupts()`. An interrupt that has been disabled by `disable_io_interrupt()` must be reenabled by `enable_io_interrupt()`.

See also: `get_nth_pci_info()`, `disable_interrupts()`

`set_isa_interrupt_handler()`, `disable_isa_interrupt()`, `enable_isa_interrupt()`

<device/KernelExport.h>

```
long set_isa_interrupt_handler(long interrupt,
                              interrupt_handler function, void *data)
```

```
long disable_isa_interrupt(long interrupt)
```

```
long enable_isa_interrupt(long interrupt)
```

These functions manage interrupts at the 8259 ISA-compatible interrupt controller. The *interrupt* is identified by its standard IRQ value.

`set_isa_interrupt_handler()` installs the handler *function* for the specified *interrupt*. This function should take one argument and return a **bool**:

```
bool handler(void *data)
```

The argument is the same *data* that's passed to `set_isa_interrupt_handler()`; it can be any kind of data the *function* might need, or `NULL`. The return value indicates whether the interrupt was handled—`TRUE` if it was and `FALSE` if not. By returning `FALSE`, the handler function can indicate that the device didn't generate the interrupt. The system can then try a different handler installed for a different device at the same interrupt number.

< However, this architecture is not currently supported, so the handler function should always return `TRUE`. >

`set_isa_interrupt_handler()` returns `B_NO_ERROR` if it can install the handler for the interrupt, and `B_ERROR` if not.

`disable_isa_interrupt()` disables the specified ISA *interrupt*, and `enable_isa_interrupt()` reenables it. Both functions return `B_ERROR` if the interrupt passed is not a valid IRQ value. Neither function considers whether the interrupt might be disabled or enabled by

some other function, such as `disable_io_interrupt()`. An interrupt that has been disabled by `disable_isa_interrupt()` must be reenabled by `enable_isa_interrupt()`.

ISA interrupts can also be managed at the Be-designed interrupt dispatcher that controls PCI interrupts. The Be interrupt controller is somewhat faster than the edge-sensitive ISA controller. If your device can generate a level-sensitive interrupt, it should use the counterpart `set_io_interrupt_handler()` function instead of `set_isa_interrupt_handler()`. However, if it depends on the edge-sensitive ISA interrupt controller widely found in the PC world, it needs to use these ISA functions.

See also: `set_io_interrupt_handler()`, `disable_interrupts()`

`spawn_kernel_thread()`

<device/KernelExport.h>

```
thread_id spawn_kernel_thread(thread_entry func, const char *name,
                              long priority, void *data)
```

This function is a counterpart to `spawn_thread()` in the Kernel Kit, which is not exported for drivers. It has the same syntax as the Kernel Kit function, but is able to spawn threads in the kernel itself.

See also: `spawn_thread()` in the Kernel Kit

`spin()`

<device/KernelExport.h>

```
void spin(double microseconds)
```

Executes a delay loop lasting at least the specified number of *microseconds*. It could last longer, due to rounding errors, interrupts, and context switches.

`start_isa_dma()`, `start_scattered_isa_dma()`, `make_isa_dma_table()`

<device/KernelExport.h>

```
long start_isa_dma(long channel, void *address, long transferCount,
                  uchar mode, uchar eMode)
```

```
long start_scattered_isa_dma(long channel, isa_dma_entry *table,
                             uchar mode, uchar eMode)
```

```
long make_isa_dma_table(void *address, long numBytes,
                       ulong numTransferBits,
                       isa_dma_entry *table, long numEntries)
```

These functions initiate ISA DMA memory transfers for the specified *channel*. They engage the ISA 8237 DMA controller.

start_isa_dma() starts the transfer of a contiguous block of physical memory beginning at the specified *address*. It requests *transferCount* number of transfers, which cannot be greater than **B_MAX_ISA_DMA_COUNT**. Each transfer will move 8 or 16 bits of memory, depending on the *mode* and *eMode* flags. These arguments correspond to the mode and extended mode flags recognized by the DMA controller.

The physical memory *address* that's passed to **start_isa_dma()** can be obtained by calling **get_memory_map()**.

start_scattered_isa_dma() starts the transfer of a memory buffer that's physically scattered in various pieces. The separate pieces of memory are described by the *table* passed as a second argument and provided by **make_isa_dma_table()**.

make_isa_dma_table() provides a description of the separate chunks of physical memory that make up the contiguous virtual buffer that begins at *address* and extends for *numBytes*. This function anticipates a subsequent call to **start_scattered_isa_dma()**, which initiates a DMA transfer. It ensures that the information it provides is in the format expected by the 8237 DMA controller. This depends in part on how many bits will be transferred at a time. The third argument, *numTransferBits*, provides this information. It can be **B_8_BIT_TRANSFER** or **B_16_BIT_TRANSFER**.

Each chunk of physical memory is described by a **isa_dma_entry** structure, which contains the following fields (not that its arcane details matter, since you don't have to do anything with the information except pass it to **start_scattered_isa_dma()**):

<code>ulong address</code>	A physical memory address (in little endian format).
<code>ushort transfer_count</code>	The number of transfers it will take to move all the physical memory at that address, minus 1 (in little endian format). This value won't be greater than B_MAX_ISA_DMA_COUNT .
<code>int flags.end_of_list:1</code>	A flag that's set to mark the last chunk of physical memory corresponding to the virtual buffer.

make_isa_dma_table() is passed a pointer to a *table* of **isa_dma_entry** structures. It fills in the table, stopping when the entire buffer of virtual memory has been described or when *numEntry* entries in the table have been written, whichever comes first. It returns the number of bytes from the virtual *address* buffer that it was able to account for in the *table*.

start_isa_dma() and **start_scattered_isa_dma()** both return **B_NO_ERROR** if successful in initiating the transfer, and **B_ERROR** if the channel isn't free.

unlock_isa_dma_channel() see **lock_isa_dma_channel()**

unlock_memory() see **lock_memory()**

write_pci_config() see **read_pci_config()**

`xpt_init()`, `xpt_ccb_alloc()`, `xpt_ccb_free()`, `xpt_action()`,
`xpt_bus_register()`, `xpt_bus_deregister()`

`<device/CAM.h>`

`long xpt_init (void)`

`CCB_HEADER *xpt_ccb_alloc(void)`

`void xpt_ccb_free(void *ccb)`

`long xpt_action(CCB_HEADER *ccbHeader)`

`long xpt_bus_register(CAM_SIM_ENTRY *entryPoints)`

`long xpt_bus_deregister(long pathID)`

These functions conform to the SCSI common access method (CAM) specification. See the draft ANSI standard *SCSI-2 Common Access Method Transport and SCSI Interface Modules* for information.

< The current implementation doesn't support asynchronous callback functions. All CAM requests are executed asynchronously in their entirety. >

Constants and Defined Types for Kernel-Loadable Drivers

This section lists the constants and types that are defined for drivers the kernel loads. Everything listed here was explained in the previous sections on “Developing a Kernel-Loadable Driver” and “Functions for Drivers”.

Constants

Control Operations

<device/Drivers.h>

Enumerated constant

B_GET_SIZE

B_SET_SIZE

B_SET_NONBLOCKING_IO

B_SET_BLOCKING_IO

B_GET_READ_STATUS

B_GET_WRITE_STATUS

B_GET_GEOMETRY

B_FORMAT

B_DEVICE_OP_CODES_END = 9999

These constants name the control operations that the kernel defines. You should expect the control hook function for any driver you develop to be called with these constants as the operation code (*op*).

All system-defined control constants are guaranteed to have values less than B_DEVICE_OP_CODES_END. Since additional constants might be defined for future releases, any that you define should be greater than B_DEVICE_OP_CODES_END.

See also: “Control Operations” on page 26

ISA DMA Transfer Maximum

<device/KernelExport.h>

<u>Defined constant</u>	<u>Value</u>
B_MAX_ISA_DMA_COUNT	0x10000

This constant indicates the maximum number of transfers for a single DMA request.

See also: `start_isa_dma()` on page 44

ISA DMA Transfer Sizes

<device/KernelExport.h>

Enumerated constant

B_8_BIT_TRANSFER
B_16_BIT_TRANSFER

These constants are passed to `make_isa_dma_table()` to indicate the size of a single DMA transfer.

See also: `start_isa_dma()` on page 44

Semaphore Flags

<device/KernelExport.h>

Defined constant

B_CAN_INTERRUPT
B_DO_NOT_RESCHEDULE
B_CHECK_PERMISSION

These flags influence the behavior of the kernel-exported functions that acquire and release semaphores.

See also: `acquire_sem_flag()` on page 33

Defined Types

cpu_status

<device/KernelExport.h>

typedef ulong **cpu_status**

This defined type is returned by **disable_interrupts()** to record whether interrupts were already disabled or not. It can be passed to **restore_interrupts()** to restore the previous state.

See also: **disable_interrupts()** on page 35

device_close_hook

<device/Drivers.h>

typedef long (***device_close_hook**)(device_info **info*)

The hook function that the kernel calls to close a device must conform to this type.

See also: “Opening and Closing a Device” on page 24

device_control_hook

<device/Drivers.h>

typedef long (***device_control_hook**)(device_info **info*, ulong *op*, void **data*)

The hook function for controlling a device must conform to this type.

See also: “Controlling the Device” on page 25

device_entry

```
<device/Drivers.h>

typedef struct {
    const char *name;
    device_open_hook open;
    device_close_hook close;
    device_control_hook control;
    device_io_hook read;
    device_io_hook write;
} device_entry
```

This structure declares the name of a device and the hook functions that the kernel can call to operate that device. The driver must provide one **device_entry** declaration for each of its devices.

See also: “Device Declarations” on page 21

device_geometry

```
<device/Drivers.h>

typedef struct {
    ulong bytes_per_sector;
    ulong sectors_per_track;
    ulong cylinder_count;
    ulong head_count;
    bool removable;
    bool read_only;
    bool write_once;
} device_geometry
```

Drivers use this structure to report the physical configuration of a mass-storage device.

See also: “B_GET_GEOMETRY” on page 27

device_info

```
<device/Drivers.h>

typedef struct {
    device_entry *entry;
    void *private_data;
} device_info
```

This structure contains publicly declared and private information about a device. It’s passed as the first argument to each of the device-specific hook functions.

See also: “Hook Functions” on page 23

device_io_hook

<device/Drivers.h>

```
typedef long (*device_io_hook)(device_info *info, void *data, ulong numBytes,  
                               ulong position)
```

The hook functions that the kernel calls to read data from or write it to a device must conform to this type.

See also: “Reading and Writing Data” on page 25

device_open_hook

<device/Drivers.h>

```
typedef long (*device_open_hook)(device_info *info, ulong flags)
```

The hook function that opens a device must conform to this type.

See also: “Opening and Closing a Device” on page 24

interrupt_handler

<device/KernelExport.h>

```
typedef bool (*interrupt_handler)(void *data)
```

The functions that are installed to handle interrupts must conform to this type.

See also: `set_io_interrupt_handler()` on page 42, `set_isa_interrupt_handler()` on page 43

isa_dma_entry

<device/KernelExport.h>

```
typedef struct {  
    ulong address;  
    ushort transfer_count;  
    uchar reserved;  
    struct {  
        int end_of_list:1;  
        int reserved:7;  
    } flags;  
} isa_dma_entry
```

This structure is filled in by `make_isa_dma_table()` and is passed unchanged to `start_scattered_isa_dma()`.

See also: `start_isa_dma()` on page 44

pci_info

<device/PCI.h>

```
typedef struct {
    ushort vendor_id;
    ushort device_id;
    uchar bus;
    uchar device;
    uchar function;
    uchar revision;
    uchar class_api;
    uchar class_sub;
    uchar class_base;
    uchar line_size;
    uchar latency;
    uchar header_type;
    uchar bist;
    uchar reserved;
    union {
        struct {
            ulong cardbus_cis;
            ushort subsystem_id;
            ushort subsystem_vendor_id;
            ulong rom_base;
            ulong rom_base_pci;
            ulong rom_size;
            ulong base_registers[6];
            ulong base_registers_pci[6];
            ulong base_register_sizes[6];
            uchar base_register_flags[6];
            uchar interrupt_line;
            uchar interrupt_pin;
            uchar min_grant;
            uchar max_latency;
        } h0;
    } u;
} pci_info
```

This structure reports values from the PCI configuration register space and describes how the device has been mapped into the system.

See also: `get_nth_pci_info()` on page 38

physical_entry

```
<device/KernelExport.h>
typedef struct {
    void *address;
    ulong size;
} physical_entry
```

This structure is used to describe a chunk of physical memory corresponding to some part of a contiguous virtual buffer.

See also: **get_memory_map()** on page 37

spinlock

```
<device/KernelExport.h>
typedef vlong spinlock
```

This data type serves the **acquire_spinlock()/release_spinlock()** protocol.

acquire_spinlock() on page 34

Developing a Driver for a Graphics Card

Like other drivers, drivers for graphics cards are dynamically loaded modules—but they’re loaded by the Application Server, the software component that’s responsible for graphics operations, not by the kernel. Because they’re add-on modules, these drivers share some similarities with other drivers:

- They lack a `main()` function, but must provide entry points where the host software (the Application Server in this case) can access driver functionality.
- They link only to their host, not to any library. They therefore can call only functions that they implement themselves or that the host explicitly exports for their benefit.
- They must be compiled as add-on modules and installed in a place well-known to the host.

However, because the host module is the Application Server, rather than the kernel, graphics card drivers must follow protocols that the Server defines, not those that the kernel imposes on other drivers. The entry points, exported functions, and installation directory are all different for graphics card drivers. Moreover, control of the driver resides only with the host Application Server; there are no functions (like `open()` or `ioctl()`) through which a program can control a graphics card driver.

If you’re developing a driver for a graphics card, disregard the preceding sections of this chapter dealing with kernel drivers and follow the rules outlined in this section instead.

Entry Points

Every graphics card driver must implement a function called `control_graphics_card()`. This is the Application Server’s main entry point into the driver; it’s the function the Server calls to set up the driver, query it for information, pass it configuration instructions, and generally control what it does. It has the following syntax:

```
long control_graphics_card(ulong op, void *data)
```

The first argument, *op*, names the operation the driver is requested to perform. The second argument, *data*, points either to some information that will help the driver carry out the request or to a location where it should write some information as a result of the operation. The exact type of data in either case depends on the nature of the operation. The *op* and *data* arguments are inextricably linked.

The return value is an error code. In general, the control function should return **B_NO_ERROR** if it can successfully respond to a particular *op* request, and **B_ERROR** if it cannot. It should also respond **B_ERROR** to any undefined *op* code requests it doesn't understand.

Control Operations

There are nine control operations (nine *op* codes defined in **device/GraphicsCard.h**) that the Application Server can request a driver's **control_graphics_card()** function to perform. These nine operations are listed and discussed in the following sections.

B_OPEN_GRAPHICS_CARD

This *op* code requests the driver to open and initialize the graphics card specified by the *data* argument. If the driver can open the card, it should do so and return **B_NO_ERROR**. If it can't, it should return **B_ERROR**. The *data* pointer refers to a **graphics_card_spec** structure with the following fields:

<code>void *screen_base</code>	The beginning of memory on the graphics card. The driver can locate the frame buffer somewhere in this memory, but not necessarily at the base address.
<code>void *io_base</code>	The base address for the I/O registers that control the graphics card. Registers are addressed by 16-bit offsets from this base address.
<code>ulong vendor_id</code>	The number that identifies the manufacturer of the graphics chip on the card.
<code>ulong device_id</code>	A number that identifies the particular graphics chip of that manufacturer.

If the driver can open the graphics card, it should take the opportunity to initialize any data structures it might need. However, it should wait for further instructions—particularly a **B_CONFIG_GRAPHICS_CARD** request—before initializing the frame buffer or turning on the video display.

If the driver returns **B_ERROR**, indicating that it's not the driver for the specified graphics card, it will immediately get a request to close the card as a prelude to being unloaded.

B_CLOSE_GRAPHICS_CARD

This operation notifies the graphics card driver that it's about to be unloaded. The *data* argument is meaningless (it doesn't point to any valid information). The Application Server ignores the return value and unloads the driver no matter what.

B_SET_INDEXED_COLOR

This operation is used to set up a color map in the driver. It requests the driver to place a particular color at a particular position in the map. The *data* argument points to a **indexed_color** structure with two pieces of information:

long index	The index of the color. This value is used as the color value in the B_COLOR_8_BIT color space.
rgb_color color	The full 32-bit color that should be associated with the index.

A driver can expect a series of **B_SET_INDEXED_COLOR** requests soon after it is opened.

B_GET_GRAPHICS_CARD_HOOKS

This *op* code requests **control_graphics_card()** to supply the Application Server with an array of function pointers. Each pointer is to a hook function that the Server can call to carry out a specific graphics task. A total of **B_HOOK_COUNT** (48 at present) pointers must be written, although only a quarter of that number are currently used. The full array should be written to the location the *data* argument points to, with **NULL** values inserted for undefined functions.

A later section, “Hook Functions” on page 60, describes the hook functions, the tasks they should perform, their arguments and return types, and their positions in the array.

B_GET_GRAPHICS_CARD_INFO

This *op* code requests the **control_graphics_card()** function to supply information about the driver and the current configuration of the screen. The *data* argument points to a **graphics_card_info** structure where it should write this information. This structure contains the following fields:

short version	The version of the Be architecture for graphics cards that the driver was designed to work with. The current version is 1; there are no earlier versions.
short id	An identifier for the driver, understood in relation to the version number. The Application Server doesn't check this number; it can be set to any value you desire.
short width	The width of the screen in pixels.
short height	The height of the screen in pixels.
short bits_per_pixel	The depth of the screen in bits per pixel. Only 32-bit (B_RGB_32_BIT) and 8-bit (B_COLOR_8_BIT) depths are currently supported.

<code>short flags</code>	A mask containing flags that describe the ability of the graphics card driver to perform particular tasks.
<code>char rgba_order[4]</code>	The order of color components as the bytes for those components are stored in video memory (in the frame buffer). This array should arrange the characters 'r' (red), 'g' (green), 'b' (blue), and 'a' (alpha) in the order in which those components are intermeshed for each pixel in the frame buffer. This field is valid only for screen depths of 32 bits per pixel.
<code>void *bits</code>	A pointer to the first byte of the frame buffer.
<code>long bytes_per_row</code>	The offset, in bytes, between two adjacent rows of pixel data in the frame buffer (the number of bytes assigned for each row).

Only two constants are currently defined for the **flags** mask:

B_CRT_CONTROL	Indicates that the driver is able to control, to any extent, the position or the size of the CRT display on the monitor—that there's a provision for controlling the CRT through software, not just hardware.
B_GAMMA_CONTROL	Indicates that the driver is able to make gamma corrections that compensate for the particular characteristics of the display device.

B_GET_REFRESH_RATES

This *op* code asks `control_graphics_card()` to place the current refresh rate, as well as the maximum and minimum rates, in the `refresh_rate_info` structure referred to by the *data* pointer. This structure contains the following fields:

<code>float min</code>	The minimum refresh rate that the graphics card is capable of, given the current configuration.
<code>float max</code>	The maximum refresh rate that the graphics card is capable of, given the current configuration.
<code>float current</code>	The current refresh rate.

All values should be provided in hertz.

B_GET_SCREEN_SPACES

This *op* code requests the driver to supply a mask containing all possible configurations of the screen space—all supported combinations of pixel depth and dimensions of the pixel grid. The mask is formed from the following constants—which are defined in **interface/InterfaceDefs.h**—and is written to the location indicated by the *data* pointer:

B_8_BIT_640x480	B_16_BIT_640x480	B_32_BIT_640x480
B_8_BIT_800x600	B_16_BIT_800x600	B_32_BIT_800x600
B_8_BIT_1024x768	B_16_BIT_1024x768	B_32_BIT_1024x768
B_8_BIT_1280x1024	B_16_BIT_1280x1024	B_32_BIT_1280x1024
B_8_BIT_1600x1200	B_16_BIT_1600x1200	B_32_BIT_1600x1200

For example, if the mask includes **B_8_BIT_1600x1200**, the driver can configure a frame buffer that's simultaneously 8 bits deep (the **B_COLOR_8_BIT** color space), 1,600 pixels wide, and 1,200 pixels high. The mask should include all configurations that the graphics card is capable of supporting.

< The Application Server currently doesn't permit depths of 16 bits. >

(**InterfaceDefs.h** defines one other screen space, **B_8_BIT_640x400**, but this is reserved for the default “supervga” driver provided by Be.)

B_CONFIG_GRAPHICS_CARD

This *op* code asks the control function to configure the display according to the values set in the **graphics_card_config** structure that the *data* argument points to. This structure contains the following fields:

ulong space	The size of the pixel grid on-screen and the depth of the frame buffer in bits per pixel. This field will be one of the constants listed above for the B_GET_SCREEN_SPACES control operation.
float refresh_rate	The refresh rate of the screen in hertz.
uchar h_position	The horizontal position of the CRT display on the monitor.
uchar v_position	The vertical position of the CRT display on the monitor.
uchar h_size	The horizontal size of the CRT display on the monitor.
uchar v_size	The vertical size of the CRT display on the monitor.

The most important configuration parameter is the **space** field. The driver should reconfigure the screen to the depth and size requested and return **B_NO_ERROR**. If it can't carry out the request, it should return **B_ERROR**.

Failure to comply with the other fields of the **graphics_card_config** structure should not result in a **B_ERROR** return value. The driver should come as close as it can to the requested refresh rate. The last four fields are appropriate only for drivers that reported that they could control the positioning of the CRT display (by setting the **B_CRT_CONTROL** flag in response to a **B_GET_GRAPHICS_CARD_INFO** request).

The values for the four CRT configuration fields range from 0 through 100, with 50 as the default. Values of less than 50 for **h_position** and **v_position** should move the display toward the left and top; those greater than 50 should move it to the right and bottom. Values of less than 50 for **h_size** and **v_size** should make the display narrower and shorter, squeezing it into a smaller area; values greater than 50 should make it wider and taller.

B_SET_SCREEN_GAMMA

This operation asks the driver to set up a table for adjusting color values to correct for the peculiarities of the display device. The *data* argument points to a **screen_gamma** structure with gamma corrections for each color component. It contains the following three fields:

<code>uchar red[256]</code>	Mappings for the red component.
<code>uchar green[256]</code>	Mappings for the green component.
<code>uchar blue[256]</code>	Mappings for the blue component.

Each field is a component-specific array. The stated color value is used as an index into the array; the value found at that index substitutes for the stated value. For example, if the value at **blue[152]** is 154, all blue component values of 152 should be replaced by 154, essentially adding to the blueness of the color as displayed.

Only drivers that indicated that could make gamma corrections (by setting the **B_GAMMA_CONTROL** flag in response to a **B_GET_GRAPHICS_CARD_INFO** request) need to respond to **B_SET_SCREEN_GAMMA** requests.

< The control function is currently not requested to perform this operation. >

Hook Functions

A graphics card driver can implement hook functions to manage the cursor and perform particular, well-defined drawing tasks on behalf of the Application Server. Drivers should implement as many of these functions as they can to speed on-screen graphics performance.

The driver informs the Application Server about these functions soon after it's loaded when its **control_graphics_card()** function receives a **B_GET_GRAPHICS_CARD_HOOKS**

request (see page 57 above). In response to this request, the driver needs to place an array of **B_HOOK_COUNT** (48) function pointers at the location the *data* argument points to.

Currently, only the first 11 slots in the array are defined. These functions fall into three groups:

- Indices 0–2: The first three functions define and manage the cursor. Drivers must implement all three of these functions, or none of them. The Application Server defers to driver-defined cursors because of the significant performance improvements they offer.
- Indices 3–9: The next seven hook functions take on specific drawing tasks, such as stroking a minimum-width line or filling a rectangle. You can choose which of these functions to implement.
- Index 10: The final function is used to synchronize the Application Server with the driver. Drivers should implement it only if the Server might sometimes need to wait for the driver to finish the drawing undertaken by any of the other hook functions.

Each undefined slot in the array of hook functions should be filled with a **NULL** pointer. Similarly, the driver should place a **NULL** value in any defined slot if it can't usefully implement the function.

Although all pointers in the array are declared to be of type **graphics_card_hook**,

```
typedef void (*graphics_card_hook)(void)
```

each function has its own set of arguments and returns a meaningful error value, declared as a **long**. The functions should be implemented to return **B_NO_ERROR** if all goes well and they're successful in performing the task at hand, and **B_ERROR** if unsuccessful. It's better by far to place a **NULL** pointer in the array than to define a function that always returns **B_ERROR**.

The coordinate system that the Application Server assumes for all hook functions equates one coordinate unit to one screen pixel. The origin is at the pixel in the left top corner of the screen. In other words, an *x* coordinate value is a left-to-right index to a pixel column and a *y* coordinate value is a top-to-bottom index to a pixel row.

The following sections discuss each of the hook functions in turn.

Index 0: Defining the Cursor

The function at index 0 is called to set the cursor image. It has the following syntax:

```
long define_cursor(uchar *xorMask, uchar *andMask, long width, long height,
                  long hotX, long hotY)
```

The first two arguments, *xorMask* and *andMask*, together define the shape of the cursor. Each mask has a depth of 1 bit per pixel, yielding a total of four possible values for each cursor pixel. They should be interpreted as follows:

<u>xorMask</u>	<u>andMask</u>	<u>meaning</u>
0	0	Transparency; let the color of the screen pixel under the cursor pixel show through.
1	0	Inversion; invert the color of the screen pixel.
0	1	White; replace the screen pixel with a white cursor pixel.
1	1	Black; replace the screen pixel with a black cursor pixel.

Inversion in its simplest form is accomplished by taking the complement of the color index or of each color component. For example:

```
color = 255 - color;
```

< However, the results of inversion may not be very pleasing given the current color map. Therefore, none of the Be-defined cursors will use inversion until the next release. It would be better for your drivers to avoid it as well. >

The second two arguments, *width* and *height*, determine the size of the cursor image in pixels. Currently, the Application Server supports only one cursor size; they must be 16 pixels wide and 16 pixels high.

The (*hotX*, *hotY*) arguments define the hot pixel in the image—the pixel that's used to report the location of the cursor. They assume a coordinate system where the pixel at the left top corner of the image is (0, 0) and the one at the right bottom corner is (15, 15).

This function should change the cursor image on-screen, if the cursor is currently displayed on-screen. But if the cursor is hidden, it should not show it. Wait for explicit calls to the next two functions to move the cursor or change its on-screen status.

Index 1: Moving the Cursor

The function at index 1 changes the location of the cursor image. It should expect two arguments:

```
long move_cursor(long screenX, long screenY)
```

In response, this function should move the cursor so that its hot pixel corresponds to (*screenX*, *screenY*).

Index 2: Showing and Hiding the Cursor

The function at index 2 shows and hides the cursor:

```
long show_cursor(bool flag)
```

If the *flag* argument is **TRUE**, this function should show the cursor image on-screen; if it's **FALSE**, it should remove the cursor from the screen.

< If this function is asked to show the cursor before the function at index 1 is called, it should show it at (0, 0). >

Index 3: Drawing a Line with an 8-Bit Color

The function at index 3 draws a straight line in the **B_COLOR_8_BIT** color space. It takes 10 arguments:

```
long draw_line_with_8_bit_depth(long startX, long startY, long endX, long endY,  
                                uchar colorIndex, bool clipToRect, short clipLeft,  
                                short clipTop, short clipRight, short clipBottom)
```

The first four arguments define the starting and ending points of the line; it begins at (*startX*, *startY*) and ends at (*endX*, *endY*). Both points are included within the line. The fifth argument, *colorIndex*, is the color of the line; it's an index into the map of 256 colors.

< In the current release, the second and third arguments are inverted; the first four arguments are ordered: *startX*, *endX*, *startY*, *endY*. >

If the sixth argument, *clipToRect*, is **TRUE**, the function should draw only the portion of the line that lies within the clipping rectangle defined by the last four arguments. The sides of the rectangle are included within the drawing area—they're inside the visible region; everything outside the rectangle is clipped.

If *clipToRect* is **FALSE**, the final four arguments should be ignored.

This function should draw a line of minimal thickness, which means a line no thicker than one pixel at any given point. If the line is more vertical than horizontal, only one pixel per row between the start and end points should be colored; if it's more horizontal than vertical, only one pixel per column should be colored.

Index 4: Drawing a Line with a 32-Bit Color

The function at index 4 is like the one at index 3, except that it draws a line in the `B_RGB_32_BIT` color space:

```
long draw_line_with_32_bit_depth(long startX, long startY, long endX, long endY,  
                                ulong color, bool clipToRect, short clipLeft,  
                                short clipTop, short clipRight, short clipBottom)
```

The only difference between this and the previous function is the *color* argument. Here the color is specified as a full 32-bit quantity with 8-bit red, green, blue, and alpha components. The *color* argument arranges the components in the order that the driver asked for them (in the `rgba_order` field of the `graphics_card_info` structure that it provided in response to a `B_GET_GRAPHICS_CARD_INFO` request).

Otherwise, this function should work just like the one at index 3.

< In the current release, the second and third arguments are inverted; the first four arguments are ordered: *startX, endX, startY, endY*. >

Index 5: Drawing a Rectangle with an 8-Bit Color

The function at index 5 should be implemented to fill a rectangle with a color specified by its index:

```
long draw_rect_with_8_bit_depth(long left, long top, long right, long bottom,  
                                uchar colorIndex)
```

The *left*, *top*, *right*, and *bottom* sides of the rectangle should be included in the area being filled.

Index 6: Drawing a Rectangle with a 32-Bit Color

The function at index 6, like the one at index 5, fills a rectangle:

```
long draw_rect_with_32_bit_depth(long left, long top, long right, long bottom,  
                                ulong color)
```

The *color* value contains the four color components—red, green, blue, and alpha—arranged in the natural order for the device (the same order that the driver recorded in the `rgba_order` field of the `graphics_card_info` structure it provided to the Application Server).

The sides of the rectangle should be included in the area being filled.

Index 7: Copying Pixel Data

The function at index 7 should copy pixel values from a source rectangle on-screen to a destination rectangle:

```
long blit(long sourceX, long sourceY, long destinationX, long destinationY,
          long width, long height)
```

The left top corner of the source rectangle is the pixel at (*sourceX*, *sourceY*). The left top pixel of the destination rectangle is at (*destinationX*, *destinationY*). Both rectangles are *width* pixels wide and *height* pixels high, and both are guaranteed to always lie entirely on-screen. The *width* and *height* arguments will always contain positive values.

Index 8: Drawing a Line Array with an 8-Bit Color

The function at index 8 should draw an array of lines in the **B_COLOR_8_BIT** color space. It takes the following set of arguments:

```
long draw_array_with_8_bit_depth(indexed_color_line *array, long numItems,
                                  bool clipToRect, short clipLeft, short clipTop,
                                  short clipRight, short clipBottom)
```

The line *array* holds a total of *numItems*. Each item is specified as an **indexed_color_line** structure, which contains the following fields:

short x1	The <i>x</i> coordinate of one end of the line.
short y1	The <i>y</i> coordinate of one end of the line.
short x2	The <i>x</i> coordinate of the other end of the line.
short y2	The <i>y</i> coordinate of the other end of the line.
uchar color	The color of the line, expressed as an index into the color map.

The function should draw each line from (**x1**, **y1**) to (**x2**, **y2**) using the **color** specified for that line.

If the *clipToRect* flag is **TRUE**, nothing should be drawn that falls outside the clipping rectangle defined by the final four arguments. The sides of the rectangle are included in the visible region. If *clipToRect* is **FALSE**, the final four arguments should be ignored.

Each line in the array should be drawn with the minimal possible thickness, as described under “Index 3: Drawing a Line with an 8-Bit Color” on page 63 above.

Index 9: Drawing a Line Array with a 32-Bit Color

The function at index 9 has the same syntax as the one at index 8, except for the first argument:

```
long draw_array_with_32_bit_depth(rgb_color_line *array, long numItems,
                                  bool clipToRect, short clipLeft, short clipTop,
                                  short clipRight, short clipBottom)
```

Here, each line in the array is specified as an **rgb_color_line** structure, rather than as an **indexed_color_line**. The two structures differ only in how the color is specified:

short x1	The <i>x</i> coordinate of one end of the line.
short y1	The <i>y</i> coordinate of one end of the line.
short x2	The <i>x</i> coordinate of the other end of the line.
short y2	The <i>y</i> coordinate of the other end of the line.
rgb_color color	The color of the line, expressed as a full 32-bit value.

In all other respects, this function should work like the one at index 8.

Index 10: Synchronizing Drawing Operations

The Application Server calls the function at index 10 to synchronize its activities with the driver. It takes no arguments:

```
long sync(void)
```

This function simply returns when the driver is finished modifying the frame buffer—when it's finished touching video RAM.

If any of the other hook functions works asynchronously—if it returns before the drawing it's asked to do is complete—the synchronizing function should wait until all drawing operations have been completed before it returns. The return value is not important; the Application Server ignores it.

However, if all the other hook functions are synchronous—if they don't return until the drawing is finished—this function would simply return; it would be empty. It's better not to implement such a function. It's preferable to put a **NULL** pointer at index 10 and save the Application Server a function call.

Therefore, your driver should implement this function only if at least one of the other hook functions draws asynchronously.

Exported Functions

Graphics card drivers are generally limited to calling functions that they implement themselves. However, they can call a handful of functions that the Application Server explicitly exports for them. The driver links to these functions in the Server, not through a library.

Functions defined in the Kernel Kit

`create_sem()`
`delete_sem()`
`acquire_sem()`
`release_sem()`

`system_time()`
`snooze()`

`spawn_thread()`
`resume_thread()`

Functions defined in the Support Kit

`atomic_add()`

Other functions

`dprintf()`
`set_dprintf_enabled()`

Functions from the first two groups are documented in the respective chapters on the Kernel Kit and the Support Kit. Functions in the last group are documented in the section on “Functions for Drivers” in this chapter on page 36.

Installation

The driver should be compiled as an add-on image, as explained in the Kernel Kit chapter. After it has been compiled, it should be installed in:

`/system/add-ons/app_server`

This is the only place where the Application Server will look for graphics card drivers to load.

The Server first looks for a driver for the graphics card in an **`app_server`** directory on a floppy disk (**`/fd/system/add-ons/app_server`**). Failing to find one, it looks next on the boot disk (**`/boot/system/add-ons/app_server`**).

When it searches each disk for a driver, the Application Server begins by looking for one developed specifically for the installed graphics card. If there are more than one, it's indeterminate which one it will choose to load. If there aren't any, the Server looks for the generic driver called **`supervga`**. This driver should be able to do a minimal job of putting

a display on-screen, but probably won't be able to exploit the full potential of the graphics card.

You can give your driver any name you wish. However, the name “supervga” is reserved for the generic driver provided by Be.

Constants and Defined Types for Graphics Card Drivers

This section lists the various constants and types that are defined for graphics card drivers. Explanations for all of them can be found in the preceding section, “Developing a Driver for a Graphics Card”.

Constants

Control Operations

<device/GraphicsCard.h>

Enumerated constant

B_OPEN_GRAPHICS_CARD
B_CLOSE_GRAPHICS_CARD
B_GET_GRAPHICS_CARD_INFO
B_GET_GRAPHICS_CARD_HOOKS
B_SET_INDEXED_COLOR
B_GET_SCREEN_SPACES
B_CONFIG_GRAPHICS_CARD
B_GET_REFRESH_RATES
B_SET_SCREEN_GAMMA

These constants define the various control operations that the Application Server can request the driver to perform.

See also: “Control Operations” on page 56

Hook Count

<device/GraphicsCard.h>

<u>Defined constant</u>	<u>Value</u>
B_HOOK_COUNT	48

This constant is the number of hook function pointers that a driver must provide. Most will be NULL pointers.

See also: “Hook Functions” on page 60

Info Flags

<device/GraphicsCard.h>

Defined constant

B_CRT_CONTROL

B_GAMMA_CONTROL

These flags report the driver's ability to control the CRT display and make gamma corrections.

See also: “**B_GET_GRAPHICS_CARD_INFO**” on page 57

Defined Types

graphics_card_config

<device/GraphicsCard.h>

```
typedef struct {  
    ulong space;  
    float refresh_rate;  
    uchar h_position;  
    uchar v_position;  
    uchar h_size;  
    uchar v_size;  
} graphics_card_config
```

This structure is used to pass the driver a set of parameters describing how the graphics card should be configured.

See also: “**B_CONFIG_GRAPHICS_CARD**” on page 59

graphics_card_hook

<device/GraphicsCard.h>

```
typedef void (*graphics_card_hook)(void)
```

This is the general type declaration for a hook function. Specific hook functions will in fact declare various sets of arguments and all return a **long** error code rather than **void**.

See also: “Hook Functions” on page 60

graphics_card_info

```
<device/GraphicsCard.h>

typedef struct {
    short version;
    short id;
    short width;
    short height;
    short bits_per_pixel;
    short flags;
    char rgba_order[4];
    void *bits;
    long bytes_per_row;
} graphics_card_info
```

Drivers use this structure to supply information about themselves and the current configuration to the Application Server.

See also: “B_GET_GRAPHICS_CARD_INFO” on page 57

graphics_card_spec

```
<device/GraphicsCard.h>

typedef struct {
    void *screen_base;
    void *io_base;
    ulong vendor_id;
    ulong device_id;
    ulong _reserved1_;
    ulong _reserved2_;
} graphics_card_spec
```

This structure informs the driver about the graphics card and how it’s mapped into the system.

See also: “B_OPEN_GRAPHICS_CARD” on page 56

indexed_color

```
<device/GraphicsCard.h>

typedef struct {
    long index;
    rgb_color color;
} indexed_color
```

This structure is used to set up the list of 256 colors in the **B_COLOR_8_BIT** color space. It locates a particular color at a particular index in the list.

See also: “**B_SET_INDEXED_COLOR**” on page 57

indexed_color_line

```
<device/GraphicsCard.h>

typedef struct {
    short x1;
    short y1;
    short x2;
    short y2;
    uchar color;
} indexed_color_line
```

This structure defines a colored line in the **B_COLOR_8_BIT** color space.

See also: “Index 8: Drawing a Line Array with an 8-Bit Color” on page 65

refresh_rate_info

```
<device/GraphicsCard.h>

typedef struct {
    float min;
    float max;
    float current;
} refresh_rate_info
```

Drivers use this structure to report the current refresh rate, and the maximum and minimum possible rates.

See also: “**B_GET_REFRESH_RATES**” on page 58

rgb_color_line

```
<device/GraphicsCard.h>
typedef struct {
    short x1;
    short y1;
    short x2;
    short y2;
    rgb_color color;
} rgb_color_line
```

This structure defines a colored line in the **B_RGB_32_BIT** color space.

See also: “Index 9: Drawing a Line Array with a 32-Bit Color” on page 66

screen_gamma

```
<device/GraphicsCard.h>
typedef struct {
    uchar red[256];
    uchar green[256];
    uchar blue[256];
} screen_gamma
```

This structure defines the table used to make gamma corrections for the screen display.

See also: “**B_SET_SCREEN_GAMMA**” on page 60

